

# A budget of paradoxes

Historical introduction to the philosophy of mathematics

András Máté

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We have proved a logical falsity from the (unlimited) comprehension using only logical rules.

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Suppose (for contradiction) that  $H_0 = f(h)$ .

$$h \in f(h) \leftrightarrow h \notin f(h)$$

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Let us substitute  $\text{v}R$  for  $x$ .

$$R(\text{v}R) \leftrightarrow \exists F(\text{v}R = \text{v}F \wedge \neg F(\text{v}R))$$

Because of the first conjunct in the scope of  $\exists$ , any concept  $F$  which makes the existential quantification true is true for just the same objects as  $R$  (because of Axiom V). Therefore, the right side is true iff  $\neg R(\text{v}R)$ .



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Let me introduce a collection of relevant paradoxes. (*A budget of paradoxes*: De Morgan 1872.)

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$$L \leftrightarrow \neg L$$

# Variants for the Liar

Liar-circle:

$$p_1 \leftrightarrow \neg p_2, p_2 \leftrightarrow \neg p_3, \dots, p_{2n-1} \leftrightarrow \neg p_1.$$

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$$L_S \leftrightarrow (L_S \text{ is not true})$$



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It is smaller than its successor.

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Known as Grelling-Nelson, Weyl, or simply heterological-paradox.

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The smallest number not definable in English by 72 characters



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$d_n = 6$  if the  $n$ th digit after the decimal point of  $a_n$  is 5 and  $d = 5$  otherwise.

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$d_n = 6$  if the  $n$ th digit after the decimal point of  $a_n$  is 5 and  $d = 5$  otherwise.

$a$  differs from any member of our sequence, but it is defined.

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The students write the test on Wednesday and they get really surprised.

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$G$  is an ordinary game between two players iff it finishes in finitely many steps.  $H$  is the following hypergame: the first player chooses an ordinary game, and then they play it. Is  $H$  an ordinary game or not?



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Russell's principle forbids self-reference. It is apparently enough to avoid the previous paradoxes.

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It is a liar-like, but infinitary paradox that does not violate the vicious circle principle and does not contain any sort of self-reference.

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Three ways out of the trap of paradoxes:

- 1 Improve logic and produce a unique general theory free of risks (logicism)
- 2 Risky theories but a reliable metatheory (formalism)
- 3 Abandon the priority of logic in favor of a more reliable basis (intuitionism)